

Analysis of Closed Networks
using
The Convolution and Mean Value Algorithms

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Analysis of a Closed Queueing Network with
Exponentially Distributed Service Times

Convolution Algorithm

Recursively calculate the
normalization constant $G(M)$
for a network with M jobs

**Mean Value Analysis
Algorithm (MVA)**

Directly calculate the performance
measures for each queue using a
recursive algorithm

Consider network of K queues with routing probabilities
 $\{p_{ij}\}$ $i, j = 1, \dots, K$ and M jobs circulating in the network

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Convolution Algorithm

- Finding the normalization constant $G(M)$ is important as it is required for obtaining the joint state probability distribution of the queueing network.

Other performance measures may then be found from the state distribution.

- For non-trivial values of K and/or M , direct computation of $G(M)$ would not be feasible.
- The *Convolution Algorithm* provides an easy, numerical approach to finding $G(M)$ recursively.

The algorithm finds $G(1), G(2) \dots G(M-1), G(M)$ in sequence

Knowing the intermediate values of $G(\cdot)$ help, as one can directly use them to compute performance measures, especially for networks of single server queues

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Some Useful Results for Closed Network of Single Server Queues

For the K queues, $j=1, \dots, K$

$$P\{n_j \geq n\} = u_j^n \frac{G(M-n)}{G(M)} \quad (5.26)$$

Actual Utilization $r_j = u_j \frac{G(M-1)}{G(M)} \quad (5.27)$

Actual Throughput $I_j = m_j r_j = m_j u_j \frac{G(M-1)}{G(M)} \quad (5.28)$

Mean Number in queue Q_j $E\{n_j\} = \sum_{m=1}^M u_j^m \frac{G(M-m)}{G(M)} \quad (5.30)$

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Convolution Algorithm for Network of Single Server Queues

$g(n,k)$ = Normalization Constant when there are n jobs circulating in a network with k queues (Q_1, \dots, Q_k) where the individual queues have the same relative utilization as in the original network

$$g(n,k) = \sum_{n_1 + \dots + n_k = n} \left(\prod_{i=1}^k u_i^{n_i} \right) \quad (5.31)$$

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Convolution Algorithm for Network of Single Server Queues (continued)

- **Initialization**

$$g(0,k) = 1 \quad g(n,1) = u_1^n \quad k=1, \dots, K \quad n=1, \dots, M$$

- **Recursion**

$$g(n,k) = g(n,k-1) + u_k g(n-1,k) \quad (5.35)$$

- **Termination**

when $g(M,K)$ has been calculated $G(M) = g(M,K)$

For $L=1, \dots, (M-1)$, the intermediate steps of the recursion give $G(L) = g(L,K)$

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**Convolution Algorithm for Network of Multi-Server Queues or
Queues with State-Dependent Service Rates**

State Dependent Service Rate at Q_i :

$m_i(j)$ = Service Rate at Q_i when it has j jobs

Multi-server Queue at Q_i (s_i servers): $m_i(j) = \min\{jm_i, s_i m_i\}$

Define $\left\{ \begin{array}{l} u_i = \frac{I_i^*}{m_i(1)} \quad i=1, \dots, K \quad (5.37) \\ A_i(n) = \prod_{j=1}^n \frac{m_i(j)}{m_i(1)} \quad i=1, \dots, K \quad (5.38) \end{array} \right.$

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**Convolution Algorithm for Network of Multi-Server Queues or
Queues with State-Dependent Service Rates**

(continued)

• **Initialization**

$$g(0, k) = 1 \quad g(n, 1) = \frac{u_1^n}{A_1(n)} \quad k=1, \dots, K \quad n=1, \dots, M$$

• **Recursion**

$$g(n, k) = \sum_{j=0}^n \frac{(u_k)^j}{A_k(j)} g(n-j, k-1) \quad (5.39)$$

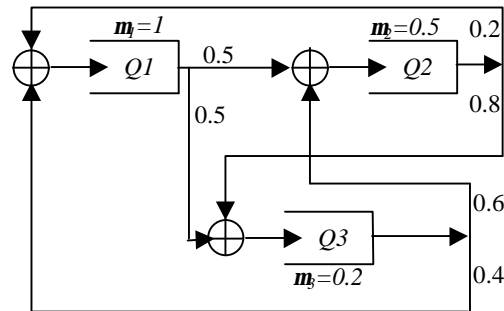
• **Termination**

when $g(M, K)$ has been calculated $G(M) = g(M, K)$

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Example



- Closed Network of Single Server, Infinite Buffer Queues
- Exponential Service Time Distribution at each queue
- Number of Circulating Jobs = $M = 4$

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Flow-Balance Equations

$$I_1^* = 0.2I_2^* + 0.4I_3^*$$

$$I_2^* = 0.5I_1^* + 0.6I_3^*$$

Choosing Q_1 as the reference queue, solve flow balance for -

Relative Throughputs $I_1^* = 1, I_2^* = 1.5385, I_3^* = 1.7308$

Relative Utilizations $u_1 = 1, u_2 = 3.077, u_3 = 8.654$

Visit Ratios $V_1 = 1, V_2 = 1.5385, V_3 = 1.7308$

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Using the Convolution Algorithm, we get -

Table 5.1. Values of $g(n,k)$ for $n=0,1,2,3,4$ and $k=1,2,3$

n	k	1	2	3
0		1	1	1
1		1	4.077	12.731
2		1	13.545	123.72
3		1	42.678	1113.35
4		1	132.32	9767.26

$$\text{Normalization Constant } G(4) = g(4,3) = 9767.26$$

$$\text{State Probability Distribution } P(n_1, n_2, n_3) = \frac{1}{9767.26} (3.077)^{n_2} (8.654)^{n_3}$$

$$\text{for } n_1, n_2, n_3 \geq 0, \text{ and } n_1 + n_2 + n_3 = 4$$

$$\text{Actual Throughputs } I_1 = 0.114, \quad I_2 = 0.1754, \quad I_3 = 0.1973$$

$$\text{Mean Number in Queue } N_1 = 0.1281, \quad N_2 = 0.5178, \quad N_3 = 3.3541$$

MVA: Mean Value Analysis Algorithm

- Directly computes performance measures of the network without computing the normalization constant or the state probability distribution
- Based on *Mean Value Theorem* (Reiser and Lavenberg)

“A customer arriving to a queue in a Product-Form Network sees the same average number in the queue as an outside observer will see if the network had one less customer”

MVA: Mean Value Analysis Algorithm

- For a network of K queues with M jobs, MVA works recursively - starting with zero jobs in the system, incrementally adding jobs until M jobs have been added to the network
- May be applied to networks of queues following any of the following service disciplines

FCFS, LCFS, PS (Processor Sharing) or IS (Infinite Number of Servers)

- Service times must be exponentially distributed

Notation

- $N_j(n)$ Mean number of jobs in Q_j when there are a total of n in the network (this includes the job currently being served at Q_j)
- $W_j(n)$ Mean time spent by a job in the queue Q_j when there are n in the network (this is the total delay at Q_j including the service time of the job)
- $1/\mathbf{m}_j$ Mean service time for a job at Q_j . Note that this will be the same regardless of the number of jobs in Q_j
- V_j Visit Ratio of Q_j . This is as defined earlier.

With Q_1 as the reference queue $V_i = \frac{I_i^*}{I_1^*} \quad i = 1, \dots, K \quad (5.18)$
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MVA Algorithm for Single Server Queues with Single Traffic Class

- **Initialization** Set $N_k(0)=0$ for each queue, i.e. $k=1, \dots, K$
- **Recursion** Do steps (1), (2) and (3) successively for $m=1, 2, \dots, M$

Step 1: Calculate for $k=1, 2, \dots, K$

$$\begin{aligned}
 W_k(m) &= \frac{1}{\mathbf{m}_k} && \text{for IS queues} \\
 &= \frac{N_k(m-1) + 1}{\mathbf{m}_k} && \text{for FCFS, LCFS, PS queues}
 \end{aligned} \tag{5.42}$$

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Step 2: Applying Little's Result

$$\text{Overall Throughput } I = \frac{m}{\sum_{k=1}^K W_k(m) V_k} \quad (5.43)$$

Step 2: Update $N_k(m)$ for $k=1, \dots, K$ as

$$N_k(m) = V_k I W_k(m) \quad (5.44)$$

- **Termination** Terminate recursion when $m=M$ is reached

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- Total average time spent by a job in the k^{th} queue (i.e. Q_k) is given directly by W_k for $k=1, \dots, K$ using (5.42) for $m=M$

(Note that this is the time spent per visit to Q_k)

- Queueing Delay for a job at Q_k for $k=1, \dots, K$ given by

$$\begin{aligned} W_{qk} &= 0 && \text{for IS queues} \\ &= W_k - \frac{1}{\mathbf{m}_k} && \text{for FCFS, LCFS, PS queues} \end{aligned} \quad (5.45)$$

- Number in Q_k given by $N_k = N_k(M)$ for $k=1, \dots, K$
- Throughput of the Network = I using (5.43) for $m=M$
- Actual Throughput of Q_k given by $I_k = I V_k$ for $k=1, \dots, K$
- Number waiting in queue in Q_k given by $N_{qk} = I_k W_{qk}$ for $k=1, \dots, K$

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MVA Algorithm for Multi-Server Queues with Single Traffic Class

(c_k servers at Q_k , where each server works with service rate \mathbf{m}_k , exponential service, FCFS or LCFS)

- **Initialization** Set $N_k(0)=0$, $p_k(0,0)=1$, and $p_k(j,0)=0$ for $j=1, \dots, (c_k-1)$ and $k=1, \dots, K$
- **Recursion** Do steps (1), (2), (3) and (4) successively for $m=1, 2, \dots, M$

Step 1: Calculate for $k=1, 2, \dots, K$

$$W_k(m) = \frac{N_k(m-1) + 1 + S_k}{\mathbf{m}_k} \quad \text{for FCFS, LCFS queues} \quad (5.47)$$

$$\text{with } S_k = \sum_{j=1}^{c_k-1} (c_k - j) p_k(j-1, m-1) \quad (5.48)$$

$(p_k(x,y))$ updated in Step 4)

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Step 2: Applying Little's Result

$$\begin{array}{l} \text{Overall} \\ \text{Throughput} \end{array} \quad \mathbf{I} = \frac{m}{\sum_{k=1}^K W_k(m) V_k} \quad (5.49)$$

Step 3: Update $N_k(m)$ for $k=1, \dots, K$ as

$$N_k(m) = V_k \mathbf{I} W_k(m) \quad (5.50)$$

Step 4: Update $p_k(j,m)$ for $k=1, \dots, K$ as

$$\begin{aligned} p_k(j, m) &= 1 - \sum_{i=1}^K p_k(i, m) && \text{for } j=0 \\ &= \frac{\mathbf{I} p_k(j-1, m-1)}{\mathbf{m}_k} && \text{for } j=1, \dots, M \end{aligned} \quad (5.51)$$

- **Termination** Terminate recursion when $m=M$ is reached

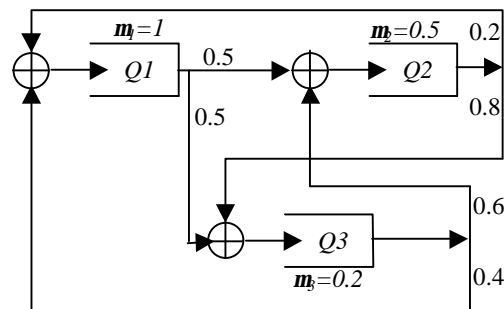
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- When recursion terminates, N_k and W_k will be directly available as $N_k=N_k(M)$ and $W_k=W_k(M)$ for $k=1, \dots, K$
- Actual network throughput I will also be available for $m=M$
- Actual throughputs of the individual queues may be found using $I_k=I V_k$ for $k=1, \dots, K$
- Average Queueing Delay at Q_k may be found as $W_{qk} = W_k - \frac{1}{m_k}$ for $k=1, \dots, K$
- Number waiting in queue in Q_k given by $N_{qk}=I_k W_{qk}$ for $k=1, \dots, K$

MVA Algorithm may also be stated for Open Networks of queues with state-dependent exponential service where there are multiple job classes. See Section 5.7.3

Example



- Closed Network of Single Server, Infinite Buffer Queues
- Exponential Service Time Distribution at each queue
- Number of Circulating Jobs = $M = 4$

Flow-Balance Equations	$I_1^* = 0.2I_2^* + 0.4I_3^*$
	$I_2^* = 0.5I_1^* + 0.6I_3^*$

Choosing Q_1 as the reference queue, solve flow balance for -

Relative Throughputs	$I_1^*=1, I_2^*=1.5385, I_3^*=1.7308$
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Relative Utilizations	$u_1=1, u_2=3.077, u_3=8.654$
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Visit Ratios	$V_1=1, V_2=1.5385, V_3=1.7308$
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- Initialize $N_1(0) = N_2(0) = N_3(0) = 0$
- Do the MVA recursion for $m = 1, 2, 3$ and 4 . (The following values were obtained at each step.)

	$W_1(1) = 1, W_2(1) = 2, W_3(1) = 5$
$m=1$	$I = 0.07855$
	$N_1(1) = 0.07855, N_2(1) = 0.2417, N_3(1) = 0.6798$

	$W_1(2) = 1.07855, W_2(2) = 2.4834, W_3(2) = 8.399$
$m=2$	$I = 0.1029$
	$N_1(2) = 0.11098, N_2(2) = 0.3932, N_3(2) = 1.49586$

$$\begin{aligned}
 m=3 \quad & W_1(3) = 1.11098, W_2(3) = 2.7864, W_3(3) = 12.4793 \\
 & \mathbf{I} = 0.1111 \\
 & N_1(3) = 0.12346, N_2(3) = 0.4764, N_3(3) = 2.4002
 \end{aligned}$$

$$\begin{aligned}
 m=M=4 \quad & W_1(4) = 1.12346, W_2(4) = 2.9528, W_3(4) = 17.001 \\
 & \mathbf{I} = 0.114 \\
 & N_1(4) = 0.12806, N_2(4) = 0.51783, N_3(4) = 3.35411
 \end{aligned}$$

Throughput of the Network = 0.114

Actual Throughputs of Q_1, Q_2, Q_3 and $Q_4 = (0.114, 0.175, 0.197)$

Actual Utilizations of Q_1, Q_2, Q_3 and $Q_4 = (0.114, 0.350, 0.985)$