

EE 679, Queuing Systems (2001-02F) Solutions to Test -6

1. Using flow balance conditions, we get the following equations.

$$\begin{aligned} I_3 &= 0.5I_1 \\ I_1 &= I + 0.2I_2 + 0.2I_3 \\ I &= 0.6I_3 + 0.5I_4 \\ I_2 &= 0.5I_1 + 0.5I_4 \end{aligned}$$

which may be solved to get

$$\begin{aligned} (I_1, I_2, I_3, I_4) &= (1.395I, 1.279I, 0.698I, 1.162I) \\ (r_1, r_2, r_3, r_4) &= (1.395r, 1.279r, 1.396r, 2.324r) \quad r = \frac{I}{m} \end{aligned}$$

(a) Network stable for $2.324r < 1 \Rightarrow I < 0.43m$

(b) In this case, $(r_1, r_2, r_3, r_4) = (0.14, 0.128, 0.14, 0.23)$

Therefore

$$P(n_1, n_2, n_3, n_4) = 0.497(0.14)^{n_1} (0.128)^{n_2} (0.14)^{n_3} (0.23)^{n_4}$$

(c) Using the result $N_i = \frac{r_i}{1 - r_i}$

we get $N_1=0.163 \quad N_2=0.147 \quad N_3=0.162 \quad N_4=0.303$

(d) $N = N_1 + N_2 + N_3 + N_4 = 0.775$

Therefore $W = \frac{N}{I} = 7.75$

2. The flow balance equations for this case are

$$\begin{aligned} I_1 &= 0.6I_2 \\ I_3 &= 0.5I_1 + 0.2I_2 + 0.4I_3 \end{aligned}$$

which gives $I_2 = 1.667I_1$ and $I_3 = 1.389I_1$

Choosing $Q1$ as the reference queue, we get the visit ratios to be

$$V_1 = 1 \quad V_2 = 1.667 \quad V_3 = 1.389$$

(a) Using the above, the steps of the MVA algorithm will be as follows.

$$\begin{array}{llll}
 (1) & m=0 & N_1=0 & N_2=0 & N_3=0 \\
 (2) & m=1 & W_1=2 & W_2=1 & W_3=2 \\
 & & \mathbf{I} = \frac{1}{2 + 1.667 + (2)(1.389)} = 0.155 & & \\
 & & N_1=0.31 & N_2=0.258 & N_3=0.431 \\
 (3) & m=2 & W_1=2.62 & W_2=1.258 & W_3=2.862 \\
 & & \mathbf{I} = \frac{2}{2.62 + (1.667)(1.258) + (1.389)(2.862)} = 0.23 & & \\
 & & N_1=0.603 & N_2=0.482 & N_3=0.914 \\
 (4) & m=3 & W_1=3.206 & W_2=1.482 & W_3=3.828 \\
 & & \mathbf{I} = \frac{3}{3.206 + (1.667)(1.482) + (1.389)(3.828)} = 0.273 & & \\
 & & N_1=0.875 & N_2=0.674 & N_3=1.451
 \end{array}$$

The mean number in each queue will be

$$N_1=0.875 \quad N_2=0.674 \quad N_3=1.451$$

(b) Note that if we choose $\mathbf{I}_1 = \mathbf{m}_1 = 0.5$, then $u_1=1$, $u_2=0.833$, $u_3=1.389$ will be the relative utilizations of the three queues. From this, it is evident that as M becomes large, i.e. $M \gg \infty$, the queue that will get bottlenecked will be Q_3 . This also implies that when M is sufficiently large, there will always be one or more users in Q_3 and hence the departure rate for this queue will approach its service rate $\mathbf{m}_3=0.5$.

Therefore, for large M , we will have

$$\mathbf{I}_3 = \mathbf{m}_3 = 0.5$$

Using this and flow balance, we get

$$\mathbf{I}_1 = 0.5/1.389 = 0.36$$

$$\text{and } \mathbf{I}_2 = (1.667)(0.36) = 0.6$$

With large M , i.e. $M \gg \infty$, we then have

$$\text{Actual Throughput} = (\mathbf{I}_1, \mathbf{I}_2, \mathbf{I}_3) = (0.36, 0.60, 0.50)$$

$$\text{Actual Utilization} = (\mathbf{r}_1, \mathbf{r}_2, \mathbf{r}_3) = (0.72, 0.60, \mathbf{r}_3 \approx 1)$$

Note that, as expected, the actual utilization of Q_3 tends toward unity as $M \gg \infty$

$$\text{Using the result } N_i = \frac{\mathbf{r}_i}{1 - \mathbf{r}_i}$$

$$\text{we get } N_1=2.57 \quad N_2=1.5 \quad N_3=(M-4.07)$$

Note that this will be how the M users will get distributed between the three queues.