

IPv6: An Introduction

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Outline

- Problems with IPv4
- Basic IPv6 Protocol
- IPv6 features
 - Auto-configuration, QoS, Security, Mobility
- Transition Plans

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Internet Protocol

Transports a datagram from source host to destination, possibly via several intermediate nodes ("routers")

Service is:

- *Unreliable*: Losses, duplicates, out-of-order delivery
- *Best effort*: Packets not discarded capriciously, delivery failure not necessarily reported
- *Connectionless*: Each packet is treated independently

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IP Datagram Header

0	4	8	16	19	31
VERS	HLEN	TOS	TOTAL LENGTH		
IDENTIFICATION		FLAG	FRAGMENT OFFSET		
TTL	PROTOCOL	CHECKSUM			
SOURCE ADDRESS					
DESTINATION ADDRESS					
OPTIONS (if any) + PADDING					

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Problems with IPv4: Limited Address Space

- IPv4 has 32 bit addresses.
- Flat addressing (only netid + hostid with "fixed" boundaries)
- Results in inefficient use of address space.
- Class B addresses are almost over.
- Addresses will exhaust in the next 5 years.
- IPv4 is victim of its own success.

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Problems with IPv4: Routing Table Explosion

- IP does not permit route aggregation (limited supernetting possible with new routers)
- Mostly only class C addresses remain
- Number of networks is increasing very fast (number of routes to be advertised goes up)
- Very high routing overhead
 - lot more memory needed for routing table
 - lot more bandwidth to pass routing information
 - lot more processing needed to compute routes

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Problems with IPv4: Header Limitations

- Maximum header length is 60 octets.
(Restricts options)
- Maximum packet length is 64K octets.
(Do we need more than that?)
- ID for fragments is 16 bits. Repeats every 65537th packet.
(Will two packets in the network have same ID?)
- Variable size header.
(Slower processing at routers.)
- No ordering of options.
(All routers need to look at all options.)

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Problems with IPv4: Other Limitations

- Lack of quality-of-service support.
 - Only an 8-bit ToS field, which is hardly used.
 - Problem for multimedia services.
- No support for security at IP layer.
- Mobility support is limited.

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IP Address Extension

- Strict monitoring of IP address assignment
- Private IP addresses for intranets
 - Only class C or a part of class C to an organization
 - Encourage use of proxy services
 - Application level proxies
 - Network Address Translation (NAT)
- Remaining class A addresses may use CIDR
- Reserved addresses may be assigned

*But these will only postpone address exhaustion.
They do not address problems like QoS, mobility, security.*

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IPng Criteria

- At least 10^9 networks, 10^{12} end-systems
- Datagram service (best effort delivery)
- Independent of physical layer technologies
- Robust (routing) in presence of failures
- Flexible topology (e.g., dual-homed nets)
- Better routing structures (e.g., aggregation)
- High performance (fast switching)
- Support for multicasting

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IPng Criteria

- Support for mobile nodes
- Support for quality-of-service
- Provide security at IP layer
- Extensible
- Auto-configuration (plug-and-play)
- Straight-forward transition plan from IPv4
- Minimal changes to upper layer protocols

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IPv6: Distinctive Features

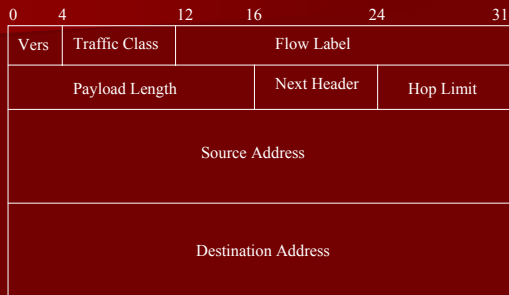
- Header format simplification
- Expanded routing and addressing capabilities
- Improved support for extensions and options
- Flow labeling (for QoS) capability
- Auto-configuration and Neighbour discovery
- Authentication and privacy capabilities
- Simple transition from IPv4

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IPv6 Header Format



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Extension Headers

- Less used functions moved to extension headers.
- Only present when needed.
- Processed only by node identified in IPv6 destination field.
=> much lower overhead than IPv4 options
Exception: Hop-by-Hop option header
- Eliminated IPv4's 40-byte limit on options
- Currently defined extension headers: Hop-by-hop, Routing, Fragment, Authentication, Privacy, End-to-end.
- Order of extension headers in a packet is defined.
- Headers are aligned on 8-byte boundaries.

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Address Types

- **Unicast** Address for a single interface.
- **Multicast** Identifier for a set of interfaces.
Packet is sent to *all* these interfaces.
- **Anycast** Identifier for a set of interfaces.
Packet is sent to the *nearest* one.

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IPv6 Addresses

- 128-bit addresses
- Multiple addresses can be assigned to an interface
- Provider-based hierarchy to be used in the beginning
- Addresses should have 64-bit interface IDs in EUI-64 format
- Following special addresses are defined :
 - IPv4-mapped
 - IPv4-compatible
 - link-local
 - site-local

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IPv6 Routing

- Hierarchical addresses are to be used.
- Initially only provider-based hierarchy will be used.
- Longest prefix match routing to be used.
(Same as IPv4 routing under CIDR.)
- OSPF, RIP, IDR, ISIS, etc., will continue as is
(except 128-bit addresses).
- Easy renumbering should be possible.
- Provider selection possible with *anycast* groups.

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QoS Capabilities

- Protocol aids QoS support, not provide it.
- Flow labels
 - To identify packets needing same quality-of-service
 - 20-bit label decided by source
 - Flow classifier: Flow label + Source/Destination addresses
 - Zero if no special requirement
 - Uniformly distributed between 1 and FFFFFFFF
- Traffic class
 - 8-bit value
 - Routers allowed to modify this field

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IPv6: Security Issues

- Provision for
 - Authentication header
 - Guarantees authenticity and integrity of data
 - Encryption header
 - Ensures confidentiality and privacy
- Encryption modes:
 - Transport mode
 - Tunnel mode
- Independent of key management algorithm.
- Security implementation is mandatory requirement in IPv6.

Mobility Support in IPv6

- Mobile computers are becoming commonplace.
- Mobile IPv6 allows a node to move from one link to another without changing the address.
- Movement can be heterogeneous, i.e., node can move from an Ethernet link to a cellular packet network.
- Mobility support in IPv6 is more efficient than mobility support in IPv4.
- There are also proposals for supporting micro-mobility.

Operation of Mobile IPv6

- Mobile node is always addressable by its **home address**.
- **Home link** is the link to which mobile nodes home address is bound.
- When attached to home link, packets are routed conventionally.
- When the node moves to foreign links, it gets a **care-of address**.
- **Binding** is an association between a home address and a care-of address.

- Care-of address is obtained using auto-configuration mechanisms of neighbour discovery.
- Mobile node, when away, registers its binding with a router on the home link called **home-agent**.
- **Binding update** and **Binding Ack** destination options are used for this purpose.
- Home agent uses **proxy neighbour discovery** to intercept packets destined for the mobile node.
- It then **tunnels** the packet to mobile node's care-of address.
- Mobile node when away uses its care-of address for communication.

- Mobile node uses a **Home Address** option to tell the other nodes its original address.
- Communicating nodes can cache the bindings and communicate with the mobile node directly.
- They use **Binding Request** destination option to learn the current binding.
- A mobile node can send a **Binding Update** to a communicating node which is using its home address as destination address.
- The communicating node should acknowledge it with a **Binding Acknowledgement**.

Neighbour Discovery

- Router Discovery - determines set of routers on the link.
- Prefix Discovery - set of on-link address prefixes.
- Parameter Discovery - to learn link parameters such as link MTU, or internet parameters like hop limit, etc.
- Address Auto-configuration - address prefixes that can be used for automatically configuring interface address.
- Address resolution - IP to link-layer address mapping.
- Duplicate Address Detection.
- Route Redirect - inform of a better first hop node to reach a particular destination.

Neighbour Discovery Operation

- Based on ICMPv6 messages
 - Router Solicitation (RS)
 - Router Advertisement (RA)
 - Neighbour Solicitation (NS)
 - Neighbour Advertisement (NA)
 - Redirect
- Router Solicitation
 - sent when an interface becomes enabled, hosts request routers to send RA immediately.

Neighbour Discovery Operation (contd..)

- Router advertisement
 - Sent by routers periodically or in response to RS.
 - Hosts build a set of default routers based on this information.
 - Provides information for address auto-configuration, set of on-link prefixes etc.
 - Supplies internet/subnet parameters, like MTU, and hop limit.
 - Includes router's link-layer address.

Neighbour Discovery Operation (contd..)

- Neighbour Solicitation
 - To request link-layer address of neighbour
 - Also used for Duplicate Address Detection
- Neighbour Advertisement
 - Sent in response to NS
 - May be sent without solicitation to announce change in link-layer address
- Redirect - used to inform hosts of a better first hop for a destination.

Address Auto-configuration

The problem

- System bootstrap ("plug and play")
- Address renumbering

Addressing Possibilities

Manual	Address configured by hand
Autonomous	Host creates address with no external interaction (e.g., link local)
Semi-autonomous	Host creates address by combining a priori information and some external information.
Stateless Server	Host queries a server, and gets an address. Server does not maintain a state.
Stateful Server	Host queries a server, and gets an address. Server maintains a state.

Auto-configuration in IPv6

- Link-local prefix concatenated with 64-bit MAC address. (*Autonomous mode*)
- Prefix advertised by router concatenated with 64-bit MAC address. (*Semi-autonomous mode.*)
- DHCPng (*for server modes*)
 - Can provide a permanent address (*stateless mode*)
 - Provide an address from a group of addresses, and keep track of this allocation (*stateful mode*)
 - Can provide additional network specific information.
 - Can register nodes in DNS.

Address Renumbering

- To migrate to a new address
 - change of provider
 - change in network architecture
- Methods
 - router adds a new prefix in RA, and informs that the old prefix is no longer valid.
 - When DHCP lease runs out, assign a new address to node.
 - DHCPng can ask nodes to release their addresses.
- Requires DNS update. DHCPng can update DNS for clients.
- Existing conversations may continue if the old address continues to be valid for some time.

Upper Layer Issues

- Minor changes in TCP
 - Maximum segment size should be based on Path MTU.
 - The packet size computation should take into account larger size of IP header(s).
 - Pseudo-header for checksum is different.
- UDP checksum computation is now mandatory.
- Most application protocol specifications are independent of TCP/IP - hence no change.
- FTP protocol exchanges IPv4 addresses - hence needs to be changed.

- The pseudo-header is changed in checksum computation:
 - Address are 128 bits.
 - Payload length is 32 bits.
 - Payload length is not copied from IPv6 header. (Extension headers should not be counted.)
 - Next header field of last extension header is used in place of protocol.
- UDP packets must also have checksum. (Since no IP checksum now.)

Changes in Other Protocols

- ICMPv6
 - Rate limiting feature added
 - Timer based
 - Bandwidth based
 - IGMP, ARP merged
 - Larger part of offending packet is included
- DNS
 - AAAA type for IPv6 addresses
 - A6 type: recursive definition of IP address
 - Queries that do additional section processing are redefined to do processing for both 'A' and 'AAAA' type records

Transition to IPv6: Design Goal

- No "flag" day.
- Incremental upgrade and deployment.
- Minimum upgrade dependencies.
- Interoperability of IPv4 and IPv6 nodes.
- Let sites transition at their own pace.
- Basic migration tools
 - Dual stack and tunneling
 - Translation

Transition Mechanisms: Dual Stack

- New nodes support both IPv4 and IPv6.
- Upgrading from IPv4 to v4/v6 does not break anything.
- Same transport layer and application above both.
- Provides complete interoperability with IPv4 nodes.

Transition Mechanism: Tunnels

- Tunnel IPv6 packets across IPv4 topology.
- Configured tunnels:
 - Explicitly configured tunnel endpoints.
 - Router to router, host to router.
- Automatic tunnels:
 - Automatic address resolution using embedded IPv4 address (like IPv4-compatible address).
 - Host to host, router to host

Transition mechanism: Translation

- This will allow communication between IPv6 only hosts and IPv4 only hosts.
- A typical translator consists of two components:
 - translation between IPv4 and IPv6 packets.
 - Address mapping between IPv4 and IPv6
- For translation, three technologies are available:
 - header conversion
 - transport relay
 - application proxy

Transition Plan for Internet

- Maintain complete V4 routing till addresses last.
- Upgrade V4 routers to dual stack.
- Incrementally build up V6 backbone routing system.
 - Use v6-over-v4 tunnels to construct 6bone.
 - Grow like Mbone (multicast backbone).
- De-activate tunnels as soon as underlying path upgraded to V6.

Transition Options for User Sites

- Incrementally upgrade V4 hosts to dual V4/V6
 - Use IPv4-compatible addresses with existing IPv4 address assignments
 - Host-to-host automatic tunneling over IPv4
- Upgrade routers to IPv6.
 - Hosts may require native IPv6 addresses
 - DNS upgrade is needed before hosts get IPv6 addresses
- Connect IPv6 router to an IPv6-enabled ISP.
- Install translators like NAT-PT or SIIT.

Thank You